Università
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## Dantzig-Wolfe: from data-driven decomposition to parallel resolution

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## Solving Prescriptive Analytics Optimization Problems



| 5 | 10.4 | 17.2 | 12 | 18.3 | 58 | 16 | 11 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |


| 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 23.3 | 0 | 15 | 17 | 102 | 51 | 0 | 237 |
| 0 | 5.6 | 0 | 0 | 0 | 0 | 0 | 0 | 12.5 |
| 0 | 0 | 0 | 0 | 0 | 6.3 | 4.3 | 2.3 | 12.5 |
| 0 | 0 | 6.2 | 0 | 6.2 | 0 | 0 | 0 | 12.5 |
| 0 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 0 | 0 | 0 | -1 | 1 | 0 | 1 |
| 7.4 | 0 | 0 | 3.1 | 0 | -2.5 | 0 | 5.3 | 15 |

## Use branch-and-cut

- problem complexity issues
- data size scalability issues


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1. to detect block structures by data driven decomposition methods
2. to exploit massive parallelism for optimization of large scale instances

## Part I: Automatic Decomposition of MIPs

Looking for decomposition schemes:

- on the problem formulation (Frangioni et al, 2010 -)
- on the MIP instance (Wang Ralphs, 2013)

Detectors search suitable decomposition patterns of a given MIP instance

- by looking for previously known structures
- by exploiting static properties of that instance (Bergner Caprara Ceselli Furini Luebbecke Malaguti Traversi, 2015)
- by communality detection (Khaniyev Elhedhli Erenay, 2018)

Generic branch-and-price-and-cut solvers:

- GCG (Gamrath Luebbecke, 2010 -)
- DIP (Ralphs Galati, 2017-)
- BapCod (Vanderbeck, 2017-)
- Coluna.jl (Marquez Nesello Vanderbeck Pessoa Sadykov, 2023)

Machine learning to choose a decomposition algorithm (Kruber
Luebbecke Parmentier, 2017)

## Promising Decompositions

Good decompositions: low running time, tight dual bound (Post-process)


## Preliminary investigation

## Research target

Investigate the link between static properties of MIP base instances and good decomposition patterns

Dataset A generation:

- 36 MIPLIB instances, generated 1000 random decompositions each
- Computed (117) static features (indep. on instance size)
- Optimized every decomposition, recorded time and bound

Experimental analysis of new decompositions from

- unknown decomp of known MIPs: base MIPs are part of Dataset A
- unknown decomp of unknown MIPs: new, unseen MIP instances

[^0]
## Mixing features: Regression for unknown problems

Independent Time and Bound regressors (XGboost). Preliminary results:
Known MIPs

Time


Time


Bound


Bound


Time prediction always possible, Bound prediction out of reach

## Feature importance

## Time (top) and bound (bottom) regressors most important features:




## A Data Driven Detection Framework



Three main components:

- D-trainer (Dataset A, xgboost regressors for Time and Bound)
- D-preprocessor
- D-optimizer (GCG)


## A Data Driven Detector

## Research target

Given a MIP instance, generate a suitable decomposition pattern


Ranking function $\mathbf{D}(\mathbf{i})$ : Percentage of decompositions dominated by decomposition i
Given decompositions i and $\mathrm{j}, i \rightarrow j$ :

$$
\begin{aligned}
& \operatorname{Time}(i) \geq \operatorname{Time}(j) \wedge \text { Bound }(i)>\operatorname{Bound}(j), \quad \text { or } \\
& \operatorname{Time}(i)>\operatorname{Time}(j) \wedge \text { Bound }(i) \geq \text { Bound }(j)
\end{aligned}
$$

[^1]
## Local Search Algorithms

## Research target

Given a base MIP instance, and a decomposition for it, improve it algorithmically


## Starting decomposition:

static detectors or data driven detector

## Neighbourhood:

set of decompositions that differ from candidate decomposition of one constraint

[^2]
## Generation

Generation: Choose one constraint from the border, insert into one block


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+ bound cannot worsen
- time prediction to avoid slow decompositions

Repeat for each constraint (border)/block combination

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Generation: Choose one constraint from the border, insert into one block


+ bound cannot worsen
- time prediction to avoid slow decompositions

Repeat for each constraint (border)/block combination
Faster computation: Sample a representative neighbourhood ( $25 x$ speedup, little loss of quality)


## Framework testing and Benchmark Configurations

Framework (D-preprocessor) configurations
D-preprocessor experimental setup

with sampling and orthogonal selection, termination: $85 \%$ convexification.

- C ++11 custom library, Boost library, Python 3.6 scripts
- Intel i7-6700K CPU and 32GB RAM
- 30 unknown MIPLIB problems
- D-optimizer: GCG 3.0.1, 5 hours timelimit
- Benchmarks: GCG, SAS-Decomp


## Root Node bounds comparison

| $\mathrm{DDW}_{\text {Sample }}$ | SAS ${ }_{\text {comm }}$ |  | GCG |  |  | GCG ${ }_{\text {Hmetis }}$ |  | GCGFull |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| run 1 | DDW Best | 10 |  | DDW Best | 6 | DDW Best | t 6 | DDW Best | 6 |
|  | Draw | 4 |  | Draw | 15 | Draw | - 14 | Draw | 14 |
|  | SAS Best | 9 |  | GCG Best | 3 | GCG Best | t 4 | GCG Best | 4 |
| run 2 | DDW Best | 10 |  | DDW Best | 7 | DDW Best | t 7 | DDW Best | 7 |
|  | Draw | 5 |  | Draw | 15 | Draw | - 14 | Draw | 14 |
|  | SAS Best | 9 |  | GCG Best | 3 | GCG Best | t 4 | GCG Best | 4 |
| run 3 | DDW Best | 10 |  | DDW Best | 8 | DDW Best | t 7 | DDW Best | 8 |
|  | Draw | 5 |  | Draw | 14 | Draw | v 13 | Draw | 13 |
|  | SAS Best | 9 |  | GCG Best | 3 | GCG Best | t 5 | GCG Best | - 4 |
|  |  |  |  |  |  |  | DDW ${ }_{\text {sample }}$ |  |  |
|  |  |  | GCG | GCG ${ }_{\text {Hmetis }}$ |  | GCG ${ }_{\text {Full }}$ | run 1 | run 2 | run 3 |
| Time [s] |  |  | 905.17 |  | 16.92 | 847.79 | 344.08 | 347.41 | 420.92 |
| Solver errors |  |  |  | 1 | 1 | 0 | 6 | 5 | 5 |
| Timeouts |  |  |  | 14 | 13 | 14 | 6 | 5 | 7 |

SAS Decomp slightly faster, 6 timeouts

## Part II: Concurrent Column

Generation

## Column Generation (CG)

## Research target

Given a base MIP instance, and a decomposition for it, solve it as quickly as possible


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## Asynchronous Column Generation (ACG)



- RMP and subproblem tasks run in parallel
- Communication through shared pool (exclusive access)
- RMP version stamps: Tasks might work with different sets of dual variables
- Generic implementation: tasks solved are solved with branch-and-cut

[^3]
## ACG example and terminating conditions



## ACG example and terminating conditions



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## ACG example and terminating conditions



## ACG example and terminating conditions



Proposition The following are sufficient and necessary conditions to guarantee convergence to the optimal solution

- each subproblem finished optimization with the latest RMP version stamp
- there are no new columns in the pool

Proof [...] $\square$

## Distributed Asynchronous Column Generation (DCG)



- RMP and Subproblem tasks run on different machines
- Communication through MPI protocol (Master-Client architecture)
- Local pool (exclusive access) on every node
- Additional handler task on every Client node


## Maximizing asynchronous computation

Experimentally, short optimization times maximize asynchronous computation:

- Dual variables are updated earlier
- RMP fetches new columns faster

Minimizing RMP optimization time:

- Rebalancing (from 10k to 20 k solutions)
- RMP update policy (store only promising solutions for the current set of dual variables)

Minimizing Subproblems optimization time:

- Short timeouts during the solve step (from 1 to 60 seconds)


## Test-bed Instances and Benchmark Algorithms

Configurations:

- CPLEX (version 12.6.3)
- Synchronous Column Generation (SCG)
- Asynchronous Column Generation (ACG)
- Distributed Asynchronous Column Generation (DCG)

Implementation: $\mathrm{C}++11$ with Boost, OpenMP and OpenMPI libraries
Experimental Setup: up to 4 machines with Intel i7-6700K CPU and 32GB RAM (Ubuntu 16.04)

Test-bed instances (root node)

- 30 Large scale instances: Multi-dimensional Variable Size and Cost Bin-packing Problem (MDVCSBP) (250, 500, 750 subproblems)
- 21 "stress" test instances: Vehicle Routing Problem with Time Windows (VRPTW) (5 subproblems)
- 21 intermediate: Multi-Depot VRPTW (155 subproblems)


## Scalability profiling

Single node (left) and multi node (right) scalability overview:


- Almost linear speed-ups for ACG on a single machine


## Scalability profiling

Single node (left) and multi node (right) scalability overview:


- Almost linear speed-ups for ACG on a single machine
- DCG scales well up to 3 machines on large scale instances
- Additional heuristic subproblems would improve performance


## Overall results

Performance profiling


30 Large scale instances (30h timeout):

- CPLEX hits 10 timeouts
- DCG is on average 87 times faster


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Performance profiling


30 Large scale instances (30h timeout):

- CPLEX hits 10 timeouts
- DCG is on average 87 times faster

21 Stress test instances (10h timeout):

- CPLEX always out of memory (19 times) or timeout (2 times)
- DCG requires on average 14 minutes


## Conclusions and Perspectives

Our data driven approach for the automatic decomposition of MIPs

- head-to-head with state-of-the-art detectors
- provides decompositions with a twisted flavor

Our parallel and distributed column generation approach

- performs one order of magnitude better than commercial solvers on large scale instances

There is still ground to cover for

- reaching commercial solvers on fully generic MIPs
- improving ML models and search algorithms
- using ML models as white boxes


[^0]:    S. Basso, A. Ceselli, A. Tettamanzi "Random Sampling and Machine Learning to Understand Good Decompositions", Annals of Operations Research 284 (2018)

[^1]:    S. Basso, A. Ceselli "Computational evaluation of ranking models in an automatic decomposition framework", Proc. of EURO/ALIO 2018 (2018)

[^2]:    S. Basso, A. Ceselli "Computational Evaluation of Data Driven Local Search for MIP Decompositions", Proc. of ODS 2019 (2019)

[^3]:    S. Basso, A. Ceselli "Asynchronous Column Generation", Proceedings of the Ninteenth Workshop on Algorithm Engineering and Experiments (ALENEX) (2017)

