



Branch-and-Price for Creating an Annual Delivery Program of Multi-Product Liquefied Natural Gas

Jørgen Glomvik Rakke^{a,b}, Guy Desaulniers^c, Henrik Andersson^a, and Marielle Christiansen^a

^a Department of Industrial Economics and Technology Management, NTNU, Trondheim, Norway

^b Norwegian Marine Technology Research Institute (MARINTEK), Trondheim, Norway

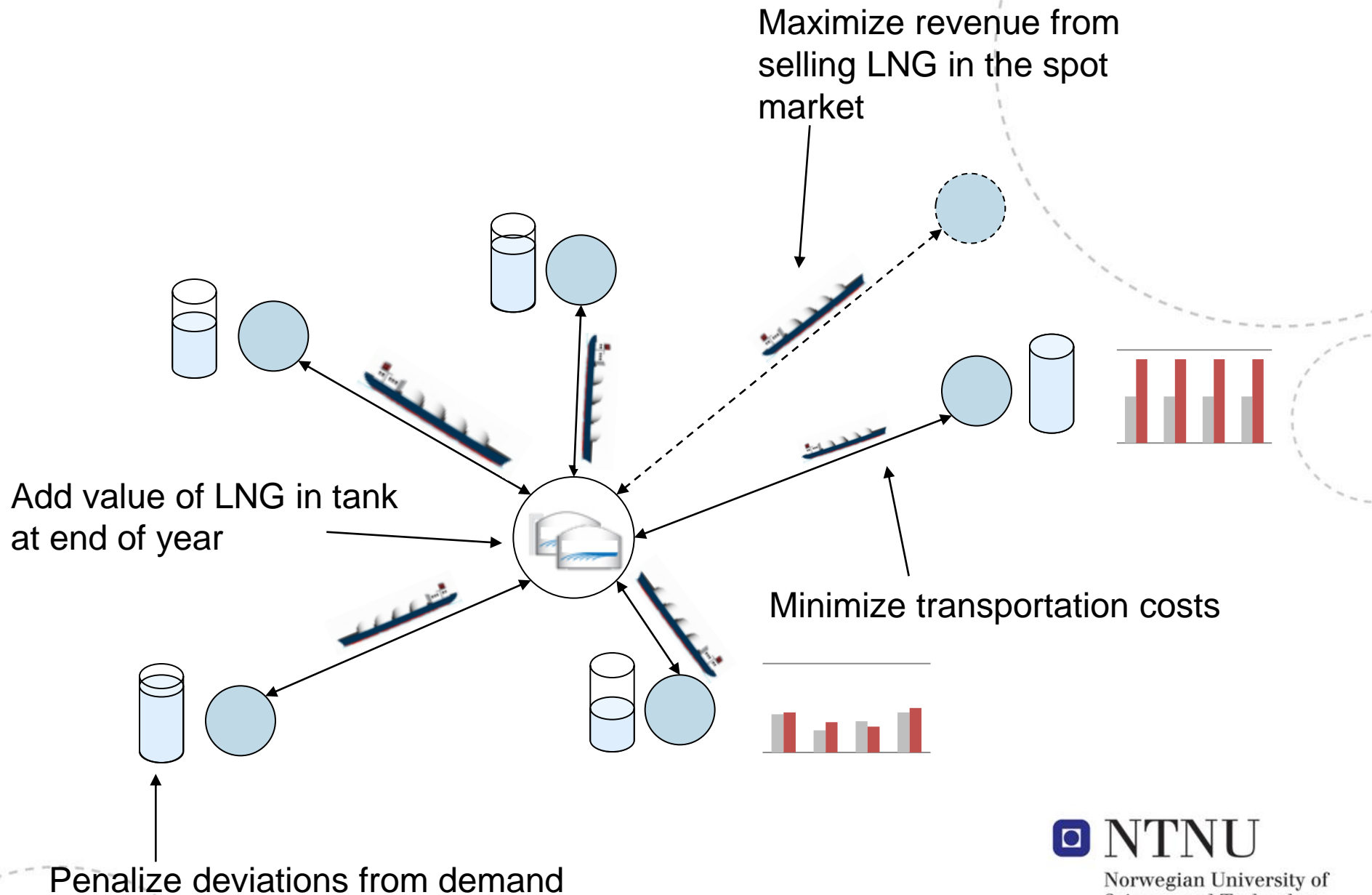
^c Département de mathématiques et génie industriel, École Polytechnique

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Outline

Problem Description
Mathematical Model
Reformulation
Subproblem
Cuts
Results



Time partitions

Full planning horizon

Partition 1

Interval 1

Interval 2

Interval 3

Partition 2

C_{c1}^{D+} C_{c1}^{D-}

Partition 1

C_{c2}^{D+} C_{c2}^{D-}

C_{c3}^{D+} C_{c3}^{D-}

C_{c4}^{D+} C_{c4}^{D-}

Partition 2

Basic Voyage Formulation (BVF)

- MIP model with scheduled voyages as binary variables
 - A scheduled voyage consists of:
 - Ship
 - Destination (Contract + terminal)
 - Departure day
- Based on pre-generation of all scheduled voyage
- The main decision is which combination of ships should deliver LNG to which contracts, and when the deliveries are to be made.

Mathematical Model

$$\begin{aligned} & \min \sum_{c \in \mathcal{C}} \sum_{v \in \mathcal{V}_c} \sum_{t \in \mathcal{T}_v} C_{cv}^T x_{cvt} \\ & + \sum_{c \in \mathcal{C}^{LT}} \sum_{i \in \mathcal{I}_c} C_{ci}^{D+} y_{ci}^+ + \sum_{c \in \mathcal{C}^{LT}} \sum_{i \in \mathcal{I}_c} C_{ci}^{D-} y_{ci}^- \\ & - \sum_{c \in \mathcal{C}^S} \sum_{v \in \mathcal{V}_c} \sum_{t \in \mathcal{T}_v} R_c^S L_v x_{cvt} - \sum_{g \in \mathcal{G}} R_g^I q_g \tau, \end{aligned}$$

(1) Objective

subject to:

$$\sum_{c \in \mathcal{C}_g \setminus \mathcal{C}^M} \sum_{v \in \mathcal{V}_c} x_{cvt} + \sum_{c \in \mathcal{C}_g \cap \mathcal{C}^M} \sum_{\substack{v \in \mathcal{V}^M \\ : t - T_{cv} - 1 \in \mathcal{T}_v}} x_{cv(t - T_{cv} + 1)} \leq B_g, \quad \forall g \in \mathcal{G}, t \in \mathcal{T},$$

(2) Berth constraints

$$q_{g,t-1} + P_{gt} - q_{gt} - \sum_{c \in \mathcal{C}_g \setminus \mathcal{C}^M} \sum_{v \in \mathcal{V}_c} L_v x_{cvt} = 0, \quad \forall g \in \mathcal{G}, t \in \mathcal{T},$$

(3)

Inventory constraints

$$\underline{Q}_{gt} \leq q_{gt} \leq \overline{Q}_{gt}, \quad \forall g \in \mathcal{G}, t \in \mathcal{T},$$

(4)

$$\sum_{c \in \mathcal{C}} \sum_{\substack{\tau \in \mathcal{T}_v \\ : t - T_{cv} < \tau \leq t}} x_{cv\tau} \leq 1, \quad \forall v \in \mathcal{V}, t \in \mathcal{T}_v,$$

(5)

Routing constraints

$$\sum_{c \in \mathcal{C}^M} \sum_{t \in \mathcal{T}_v^M} x_{cvt} = 1, \quad \forall v \in \mathcal{V}^M,$$

(6)

Maintenance constraints

$$\sum_{v \in \mathcal{V}_c} \sum_{t \in \mathcal{T}_i^I \cap \mathcal{T}_v} L_v x_{cvt} + y_{ci}^- - y_{ci}^+ = D_{ci}, \quad \forall c \in \mathcal{C}^{LT}, i \in \mathcal{I}_c,$$

(7)

Contractual constraints

$$x_{cvt} \in \{0, 1\}, \quad \forall c \in \mathcal{C}, v \in \mathcal{V}_c, t \in \mathcal{T}_v,$$

(8)

$$y_{ci}^+, y_{ci}^- \geq 0, \quad \forall c \in \mathcal{C}^{LT}, i \in \mathcal{I}_c.$$

(9)

Objective function

$$\begin{aligned}
 & \min \sum_{c \in \mathcal{C}} \sum_{v \in \mathcal{V}_c} \sum_{t \in \mathcal{T}_v} C_{cv}^T x_{cvt} && \text{Transportation cost} \\
 & + \sum_{c \in \mathcal{C}^{LT}} \sum_{i \in \mathcal{I}_c} C_{ci}^{D+} y_{ci}^+ + \sum_{c \in \mathcal{C}^{LT}} \sum_{i \in \mathcal{I}_c} C_{ci}^{D-} y_{ci}^- && \text{Penalty costs} \\
 & - \sum_{c \in \mathcal{C}^S} \sum_{v \in \mathcal{V}_c} \sum_{t \in \mathcal{T}_v} R_c^S L_v x_{cvt} - \sum_{g \in \mathcal{G}} R_g^I q_{g\bar{T}}, && \text{Spot revenue}
 \end{aligned}$$

Problems

- Weak LP-relaxation
 - Mainly introduced by the penalties for over and under delivery

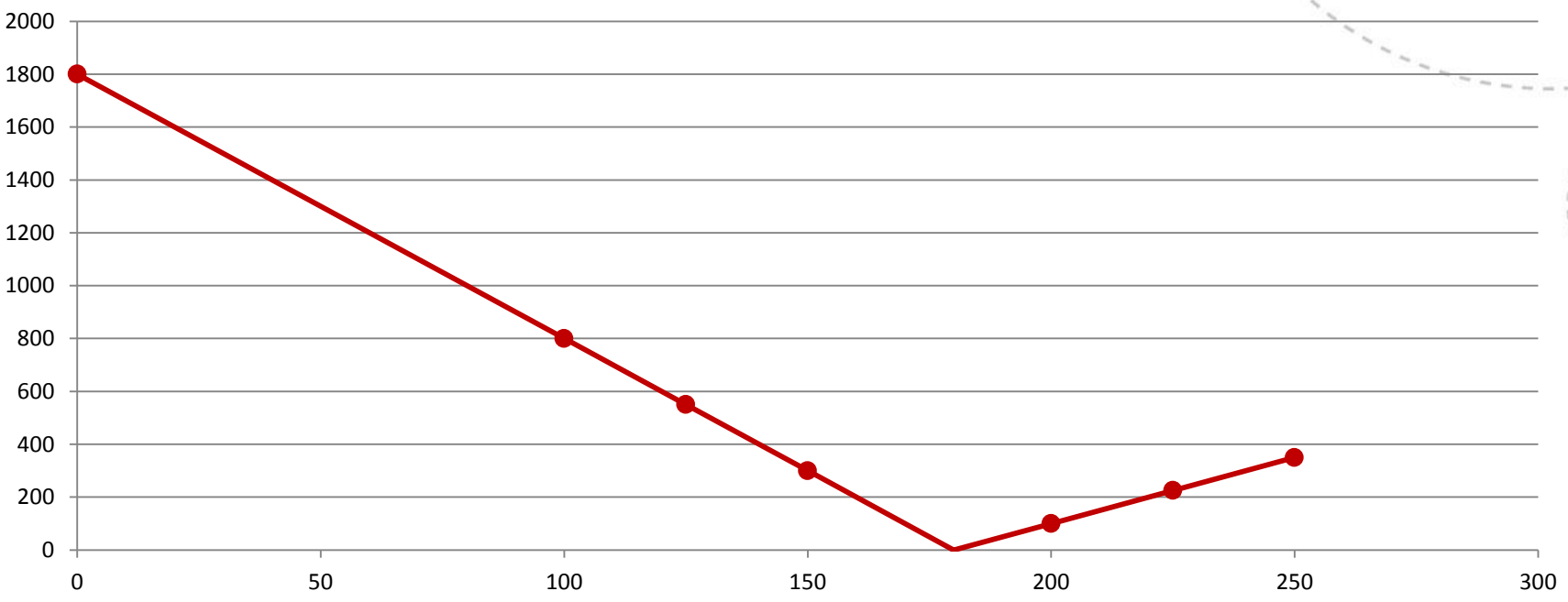
$$\begin{aligned}
 & \min \sum_{c \in \mathcal{C}} \sum_{v \in \mathcal{V}_c} \sum_{t \in \mathcal{T}_v} C_{cv}^T x_{cvt} \\
 & + \sum_{c \in \mathcal{C}^{LT}} \sum_{i \in \mathcal{I}_c} C_{ci}^{D+} y_{ci}^+ + \sum_{c \in \mathcal{C}^{LT}} \sum_{i \in \mathcal{I}_c} C_{ci}^{D-} y_{ci}^- \\
 & - \sum_{c \in \mathcal{C}^S} \sum_{v \in \mathcal{V}_c} \sum_{t \in \mathcal{T}_v} R_c^S L_v x_{cvt} - \sum_{g \in \mathcal{G}} R_g^I q_g \bar{T},
 \end{aligned}$$

Example

- Heterogeneous fleet, ship capacities = 100, 125, 150
- Demand = 180
- Under-delivery penalty = 10
- Over-delivery penalty = 5
- Minimum under-delivery = $180 - 150 = 30$
- Minimum over-delivery = $100 + 100 - 180 = 20$

Penalty function

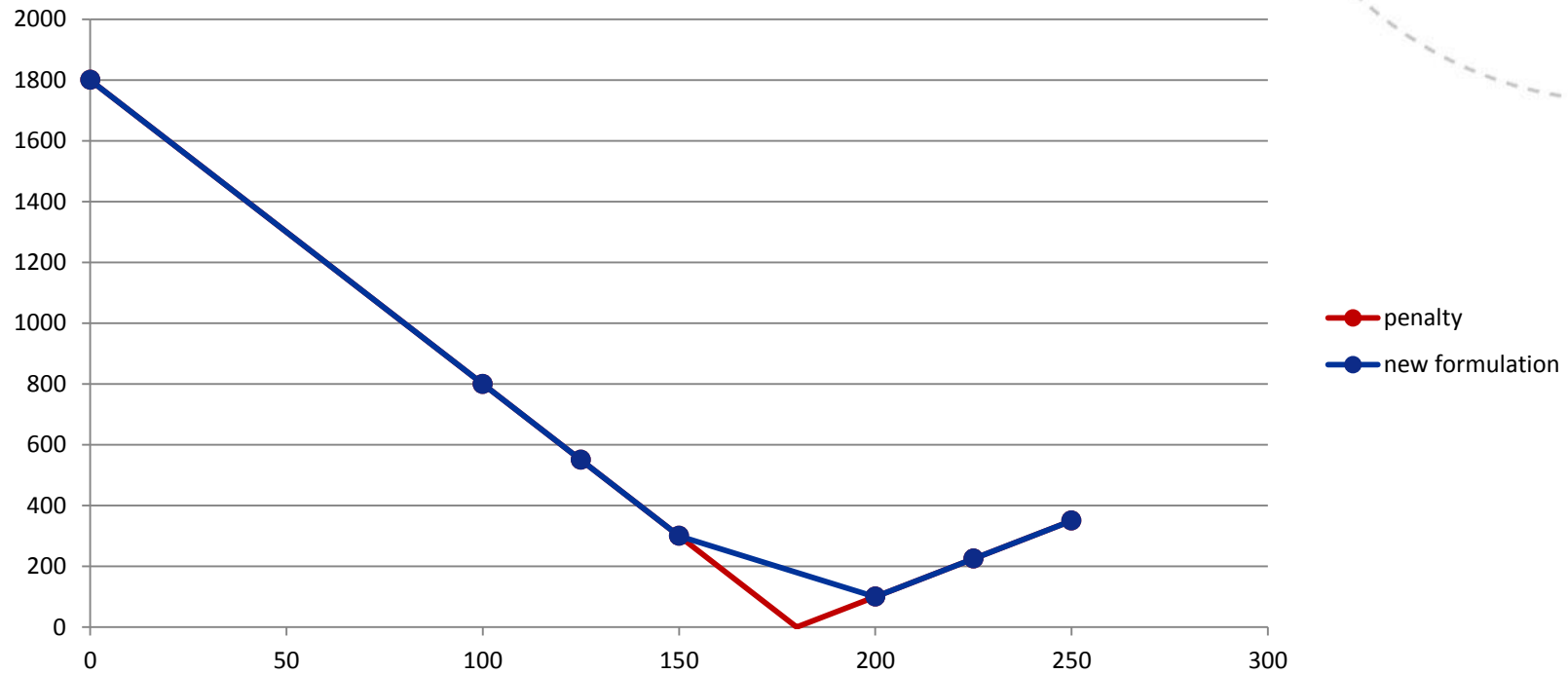
penalty



How to deal with this problem

- Create delivery patterns
 - A delivery pattern specifies:
 - Contract
 - Time period
 - Number of deliveries for all ships

Penalty function



Pattern Based Formulation (PBF)

- Set of delivery patterns for each time period and contract \mathcal{P}_{ci}
- Pattern variable z_{cip}

$$\begin{aligned} \min \quad & \sum_{c \in \mathcal{C}} \sum_{v \in \mathcal{V}_c} \sum_{t \in \mathcal{T}_v} C_{cv}^T x_{cvt} + \sum_{c \in \mathcal{C}^{LT}} \sum_{i \in \mathcal{I}_c} \sum_{p \in \mathcal{P}_{ci}} C_{cip}^P z_{cip} \\ & - \sum_{c \in \mathcal{C}^S} \sum_{v \in \mathcal{V}_c} \sum_{t \in \mathcal{T}_v} R_c^S L_v x_{cvt} - \sum_{g \in \mathcal{G}} R_g^I q_g \bar{T}, \end{aligned} \quad (10)$$

subject to:

constraints (2)–(6),

$$\sum_{t \in \mathcal{T}_i \cap \mathcal{T}_v^I} x_{cvt} - \sum_{i \in \mathcal{P}_{ci}} N_{cipv} z_{cip} = 0, \quad \forall c \in \mathcal{C}^{LT}, i \in \mathcal{I}_c, v \in \mathcal{V}_c, \quad (11) \quad \text{Linking constraints}$$

$$\sum_{i \in \mathcal{P}_{ci}} z_{cip} = 1, \quad \forall c \in \mathcal{C}^{LT}, i \in \mathcal{I}_c, \quad (12) \quad \text{Convexity constraints}$$

$$x_{cvt} \in \{0, 1\}, \quad \forall c \in \mathcal{C}, v \in \mathcal{V}_c, t \in \mathcal{T}_v, \quad (13)$$

$$z_{cip} \geq 0, \quad \forall c \in \mathcal{C}^{LT}, i \in \mathcal{I}_c, p \in \mathcal{P}_{ci}. \quad (14)$$

How to generate delivery patterns

- Pre generate all possible delivery patterns
 - Only feasible for really small test instances
- Column generation
 - Split the problem into a master and sub problems to generate delivery patterns with negative reduced cost

Subproblem

$$\overline{C}_{cip} = C_{cip}^P - \alpha_{ci} + \sum_{v \in \mathcal{V}_c} N_{cipv} \beta_{civ}$$

$$\min \quad C^{D+} y^+ + C^{D-} y^- - \alpha + \sum_{v \in \mathcal{V}_c} \beta_v n_v$$

$$\text{s.t.} \quad \sum_{v \in \mathcal{V}_c} L_v n_v + y^- - y^+ = D,$$

$$l_v \leq n_v \leq u_v, \text{ integer, } \forall v \in \mathcal{V}_c,$$

$$y^+ \geq 0,$$

$$y^- \geq 0,$$

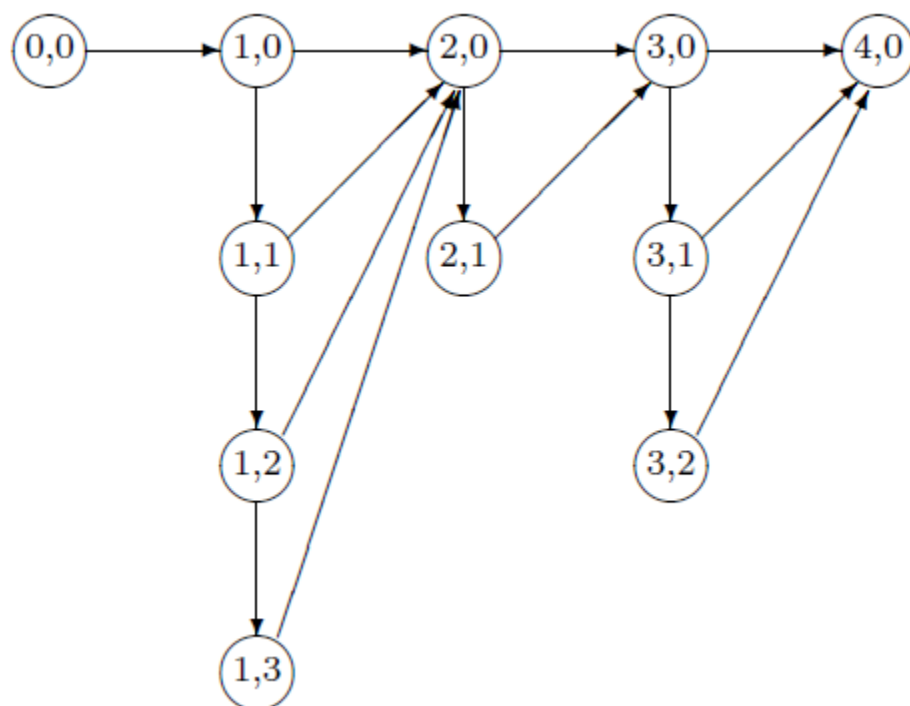
Subproblem

- If $\beta_v - C^{D-}L_v \geq 0$, n_v can be set to l_v and removed
- If $\beta_v + C^{D+}L_v < 0$, n_v can be set to u_v and removed
- Let $\bar{\mathcal{V}}_c \subseteq \mathcal{V}_c$ be the set of vessels such that $\beta_v < -C^{D+}L_v$
- Define \tilde{D} as $D - \sum_{v \in \mathcal{V}_c} L_v l_v - \sum_{v \in \bar{\mathcal{V}}_c} L_v (u_v - l_v)$
- Let $\mathcal{V}_c^- \subseteq \mathcal{V}_c$ be the set of vessels such that $-C^{D+}L_v < \beta_v < C^{D-}L_v$
- Let $u'_v = u_v - l_v$

Subproblem

$$\begin{aligned}
 \min \quad & K + (C^{D+} + C^{D-})y^+ + \sum_{v \in \mathcal{V}_c^-} (\beta_v - C^{D-}L_v)n_v \\
 \text{s.t.} \quad & \sum_{v \in \mathcal{V}_c^-} L_v n_v \leq \tilde{D} + y^+, \\
 & 0 \leq n_v \leq u'_v, \text{ integer, } \forall v \in \mathcal{V}_c^-, \\
 & y^+ \geq 0.
 \end{aligned}$$

$$K = C^{D-}D - \alpha + \sum_{v \in \mathcal{V}_c} (\beta_v - C^{D-}L_v)l_v + \sum_{v \in \bar{\mathcal{V}}_c} (\beta_v - C^{D-}L_v)u'_v$$



$$c_{ij} = \begin{cases} K & \text{if } i = (0,0) \\ k(\beta_v - C^D - L_v) & \text{if } (i,j) \text{ is a select arc and } i = (v,k) \\ 0 & \text{otherwise.} \end{cases}$$

$$q_{ij} = \begin{cases} kL_v & \text{if } (i,j) \text{ is a select arc and } i = (v,k) \\ 0 & \text{otherwise.} \end{cases}$$

Labels

- A label contains two elements
 - The accumulated cost Z
 - The accumulated quantity Q
- By using arc (i, j) you update label $E_i = (Z_i, Q_i)$ to label $E_j = (Z_j, Q_j)$

$$Q_j = Q_i + q_{ij}$$

$$Z_j = Z_i + c_{ij} + z_{ij}(Q_i)$$

$$z_{ij}(Q_i) = \begin{cases} 0 & \text{if } Q_i + q_{ij} \leq \tilde{D} \\ \infty & \text{if } q_{ij} > 0 \text{ and } Q_i \geq \tilde{D} \\ (C^{D+} + C^{D-})(Q_i + q_{ij} - \tilde{D}) & \text{otherwise.} \end{cases}$$

Dominance

- Consider two labels

$$E_1 = (Z_1, Q_1)$$

$$E_2 = (Z_2, Q_2)$$

- E_1 dominates E_2 if:

$$Z_1 \leq Z_2 \text{ and } Q_1 \leq Q_2$$

- Given that the amount delivered is not constrained we can strengthen this dominance rule to:

$$Z_1 + (C^{D+} + C^{D-}) \max\{0, Q_1 - Q_2\} \leq Z_2$$

- Where we can discard label E_2 even if $Q_2 < Q_1$

Cuts

- Timing cuts (similar to the ones presented in Engineer et al. (2012))
- Minimum number of loadings
- Maximum number of loadings

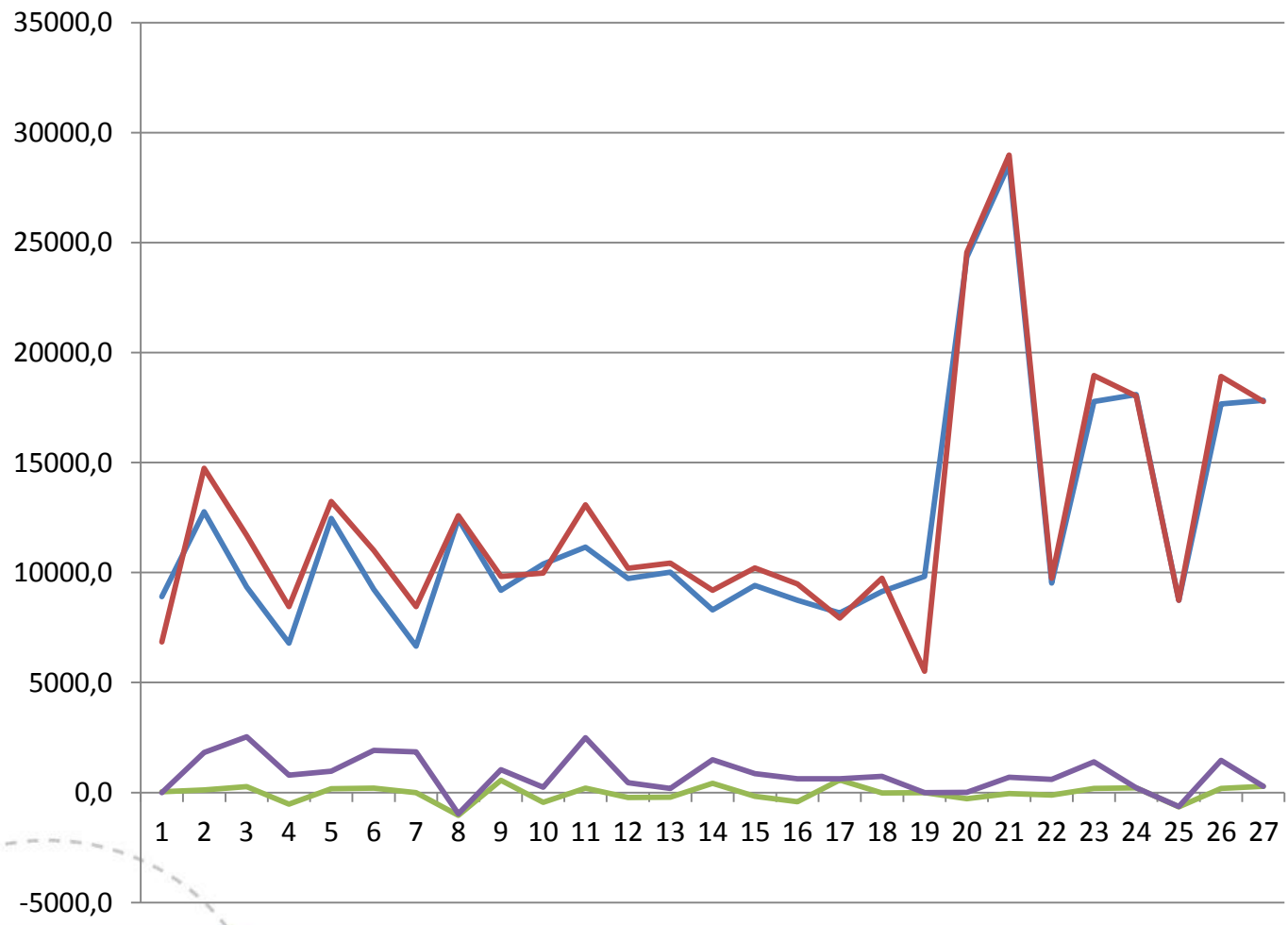
Solution method

- Branch-price-and-cut
 - Subproblem solved as a shortest path problem with resource constraints
 - The n best patterns priced in until no negative reduced cost pattern found
 - Generate only the k most violated cuts, until no violated cuts found

Tests

- Root node solutions of BVF vs. PBF
- Effect of cuts
- BVF vs. PBF

LB comparison PBF vs. BVF

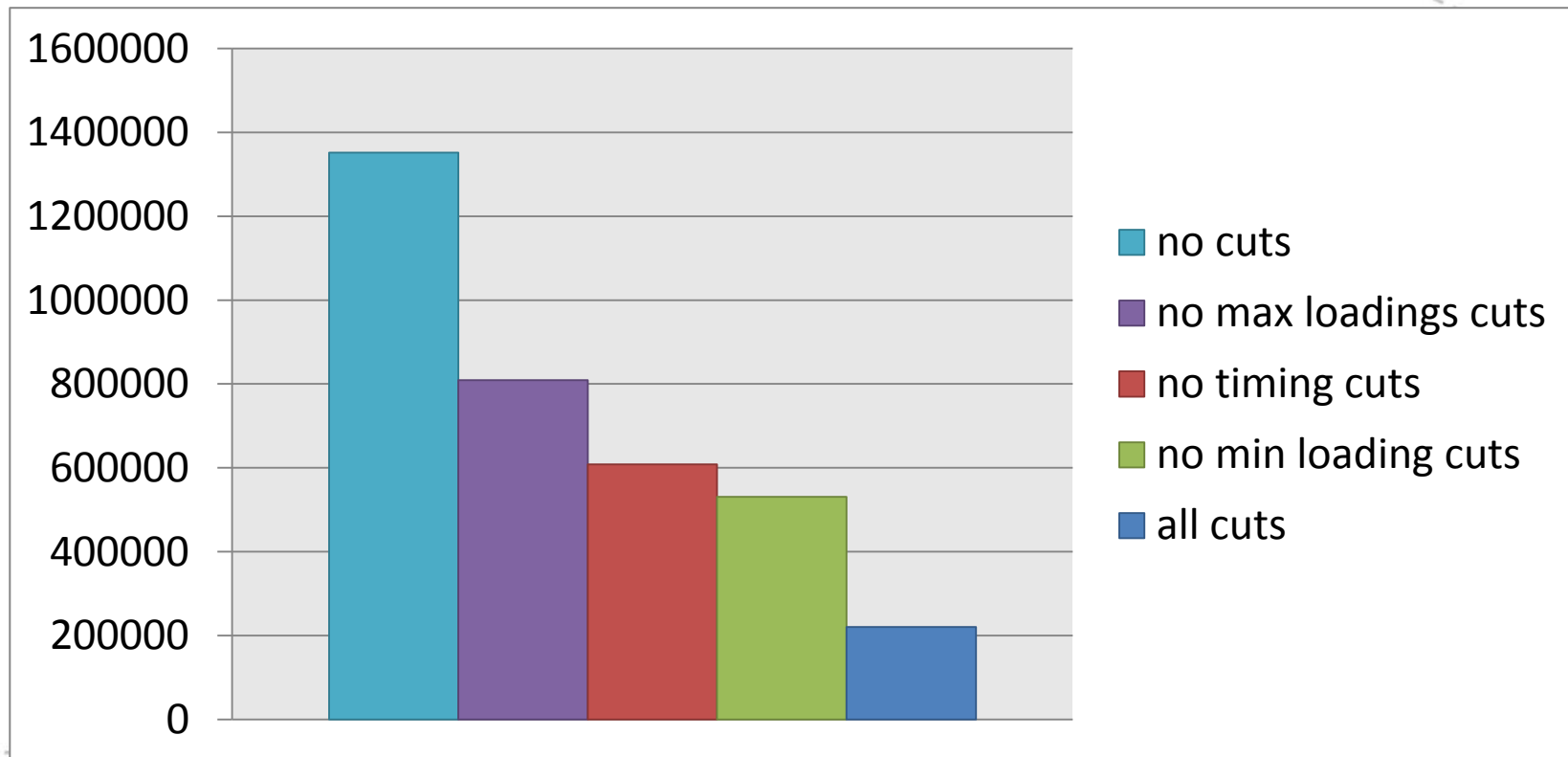


Difference =
PBF-BVF

- root diff (no cuts)
- diff (no cuts)
- root diff (BVF with cuts)
- diff (BVF with cuts)

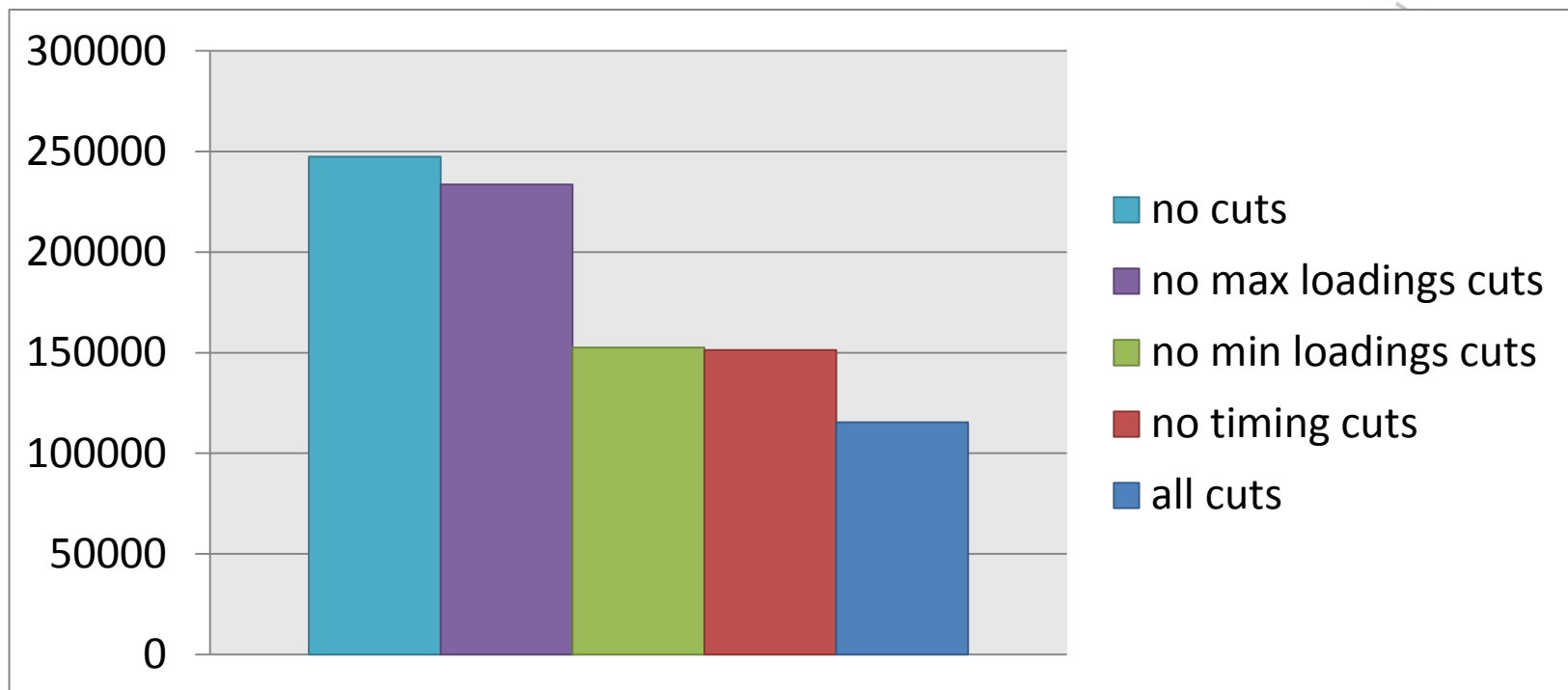
Effect of cuts for PBF

Number of nodes



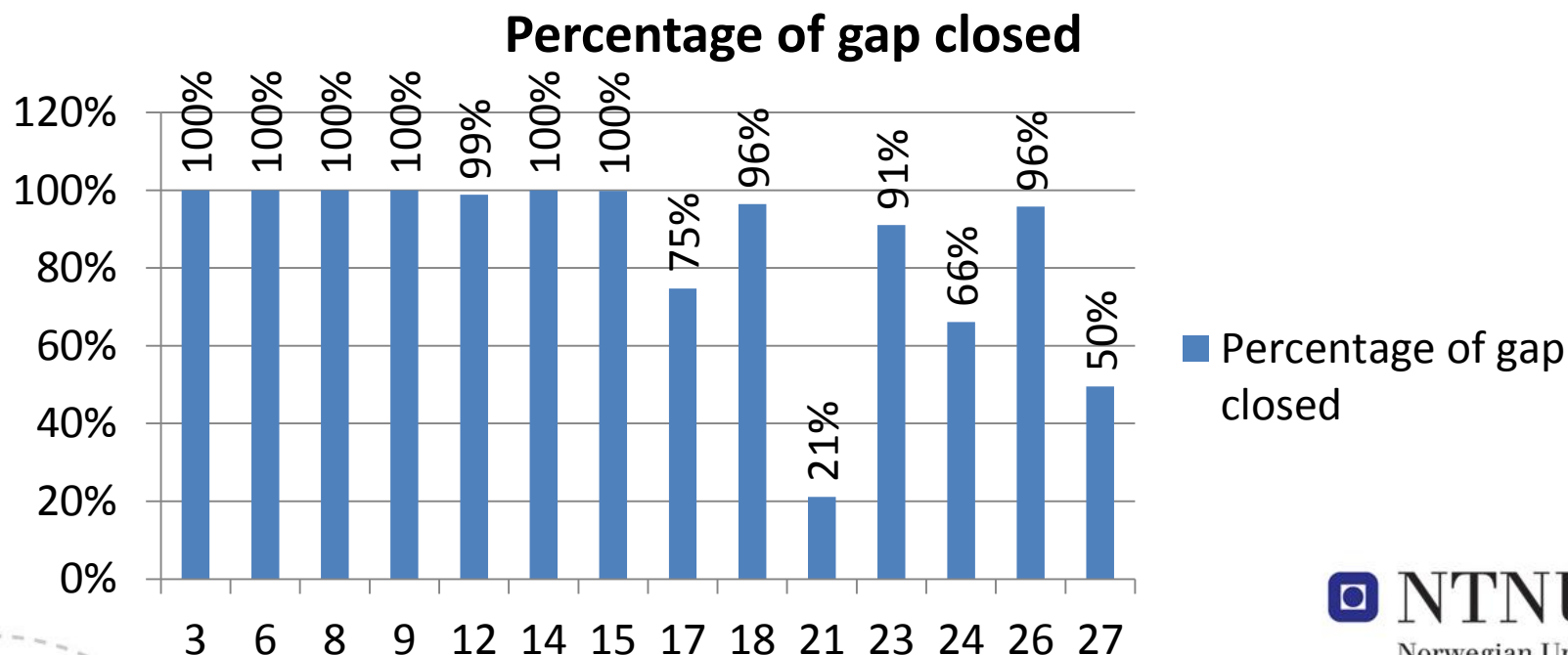
Effect of cuts for PBF

Solution time



Gap closed by the PBF

- 13 instances solved by both solution methods
 - Instances 1,2,4,5,7,10,11,13,16,19,20,22 and 25



Concluding remarks

- Results show that the new formulation has a much stronger LP relaxation
- Solve more instances to optimality
- Close most of the gap for the remaining instances



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^a Department of Industrial Economics and Technology Management, NTNU, Trondheim, Norway

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