The Mesh Adaptive Direct Search Algorithm for Discrete Blackbox Optimization

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Presentation outline

Blackbox optimization

Motivating example

The MADS algorithm

Computational experiments

Discussion

Blackbox optimization (BBO) problems

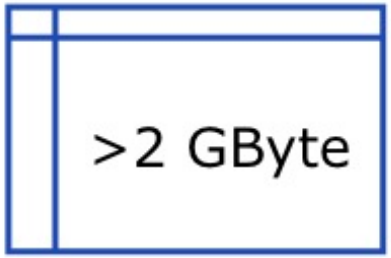
- Optimization problem:

$$\min_{x \in \Omega} \quad f(x)$$

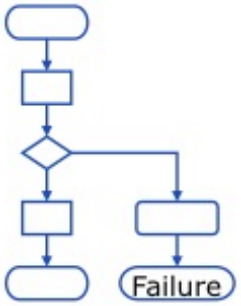
- Evaluations of f (the objective function) and of the functions defining Ω are usually the result of a computer code (a blackbox).
- Variables are typically continuous, but in this work, some of them are discrete – integers or *granular* variables.

Blackboxes as illustrated by J. Simonis [ISMP 2009]

Long runtime

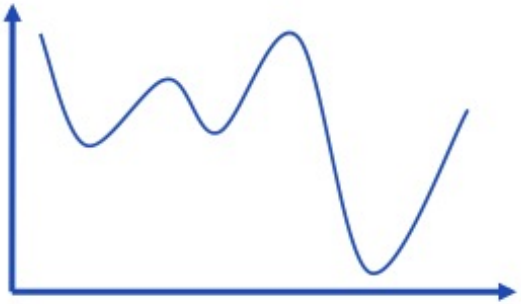


Large memory requirement



Software might fail

No derivatives available



Local optima

Non-smooth, noisy

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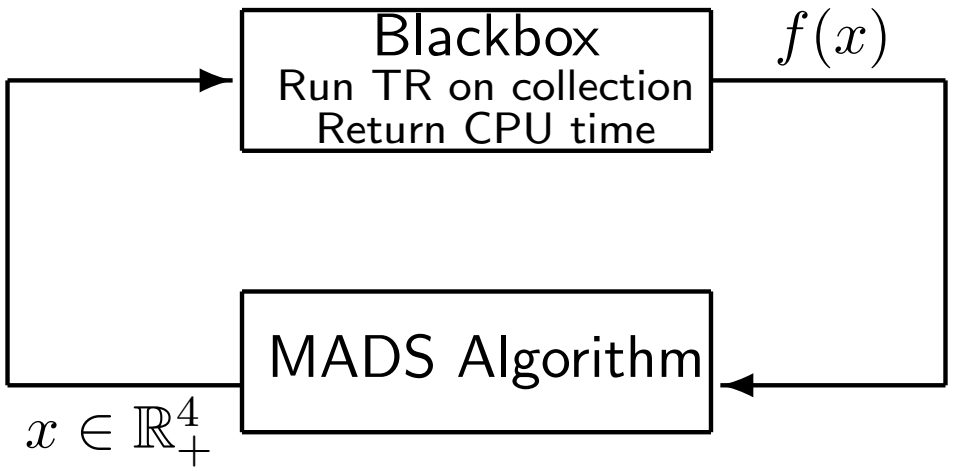
Motivating example: Parameters tuning (1/2)

- [Audet and Orban, 2006].
- The classical Trust-Region algorithm depends on four parameters $x = (\eta_1, \eta_2, \alpha_1, \alpha_2) \in \mathbb{R}^4_+$.
- Consider a collection of 55 test problems from `CUTEr`.
- Let $f(x)$ be the CPU time required to solve the collection of problems by a Trust-Region algorithm with parameters x.
- $f(x_0) \simeq 3h45$ with the textbook values $x_0 = (1/4,\ 3/4,\ 1/2,\ 2)$.

Motivating example: Parameters tuning (2/2)



- This optimization produced $\hat{x}$ with $f(\hat{x}) \simeq 2h50$.

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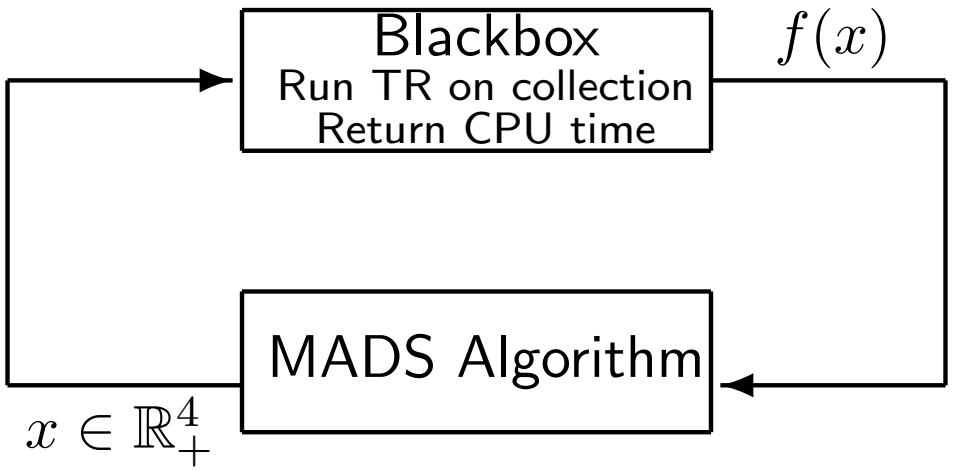

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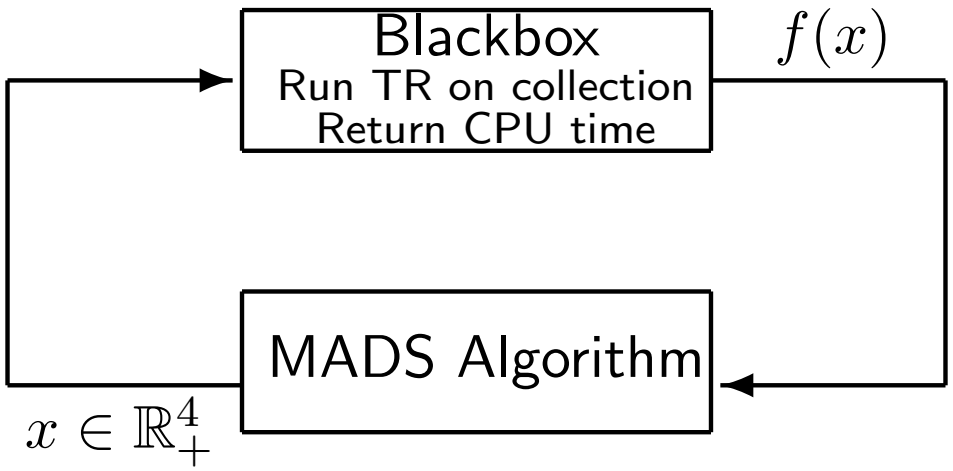

- This optimization produced $\hat{x}$ with $f(\hat{x}) \simeq 2h50$.Victory ?
  No because
  $\hat{x} = (0.22125,\ 0.94457031,\ 0.37933594,\ 2.3042969)$.

Granular variables

The initial point $x_0 = (0.25,\ 0.75,\ 0.50,\ 2.00)$ is frequently used because each entry is a multiple of 0.25.

Its granularity is $\mathcal{G} = 0.25$

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- How can we devise a direct search algorithm so that it stops on a prescribed granularity?
  With a granularity of $\mathcal{G} = 0.05$, the code might produce
  $$\hat{x} = (0.20, 0.95, 0.40, 2.30).$$
  Which is much nicer (for a human) than
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  Which is much nicer (for a human) than
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- This may be achieved using integer variables, together with a relative scaling, but there is a simpler way for mesh-based methods.

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Mesh Adaptive Direct Search (MADS) in $\mathbb{R}^n$

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- One iteration = search and poll.
- The search allows trial points generated anywhere on the mesh.
- The poll consists in generating a list of trial points constructed from poll directions. These directions grow dense.
- At the end of the iteration, the mesh size is reduced if no new success point is found.

[0] Initializations (x_0, Δ_0: initial poll size)

[1] Iteration k

- let $\delta^k \leq \Delta^k$ be the mesh size parameter
- **Search**
  - test a finite number of mesh points
- **Poll** (if the Search failed)
  - construct set of directions D_k
  - test poll set $P_k = \{x_k + \delta^k d : d \in D_k\}$ with $\|\delta^k d\| \simeq \Delta_k$

[2] Updates

- if success
  - $x_{k+1} \leftarrow$ success point
  - increase Δ^k
- else
  - $x_{k+1} \leftarrow x_k$
  - decrease Δ^k
- $k \leftarrow k+1$, stop if $\Delta^k \leq \Delta_{\min}$ or go to **[1]**

Poll illustration (successive fails and mesh shrinks)

$$\delta^k = 1$$
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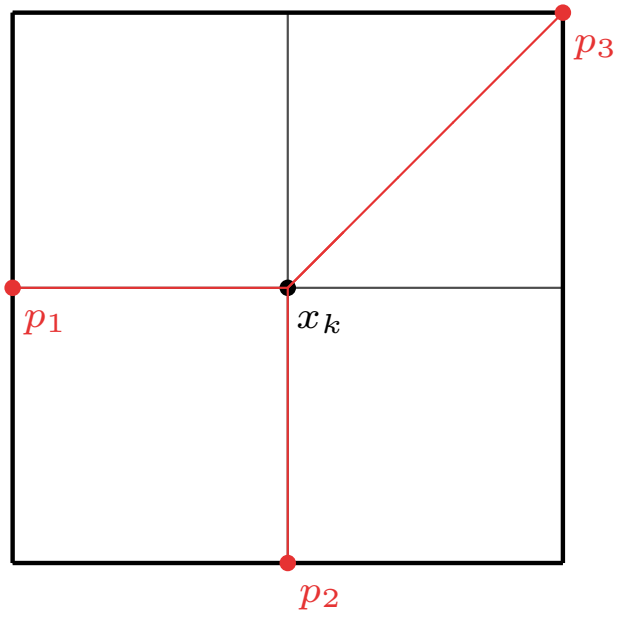

trial points$=\{p_1, p_2, p_3\}$

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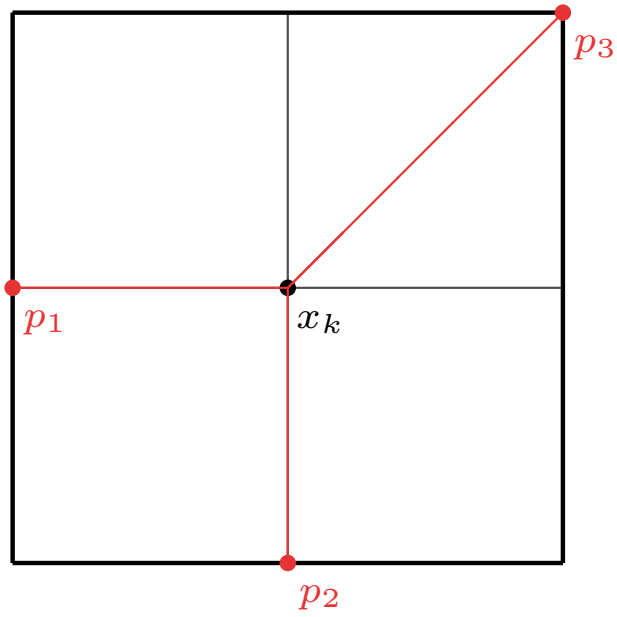
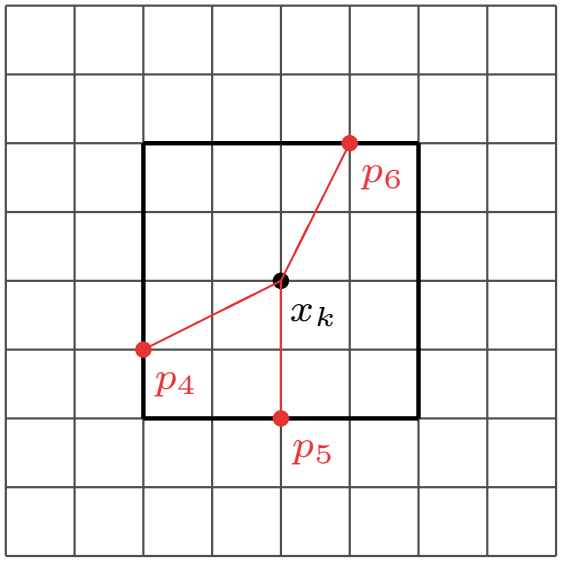
$$\delta^k = 1 \qquad \Delta^k = 1$$

$$\delta^{k+1} = 1/4 \qquad \Delta^{k+1} = 1/2$$





trial points$=\{p_1, p_2, p_3\}$ $\qquad = \{p_4, p_5, p_6\}$

Poll illustration (successive fails and mesh shrinks)

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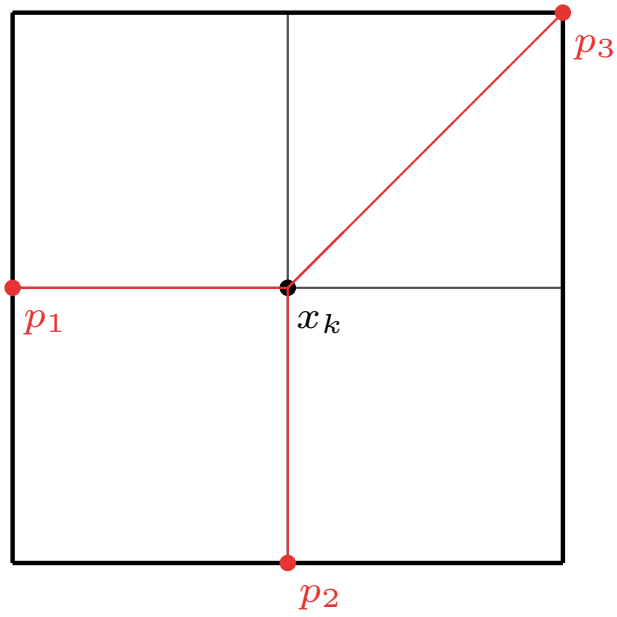

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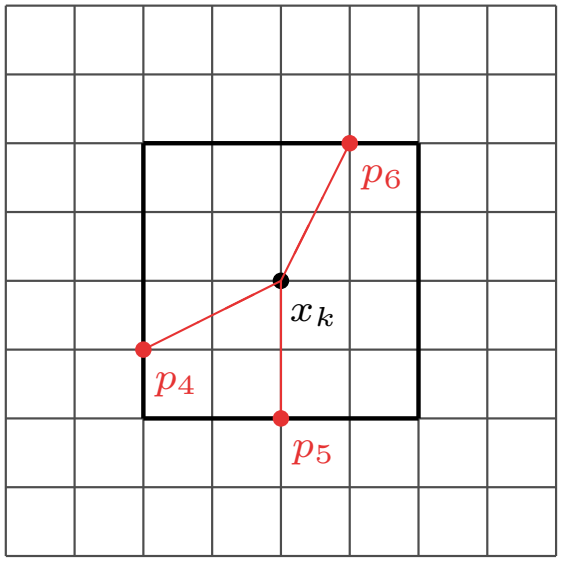
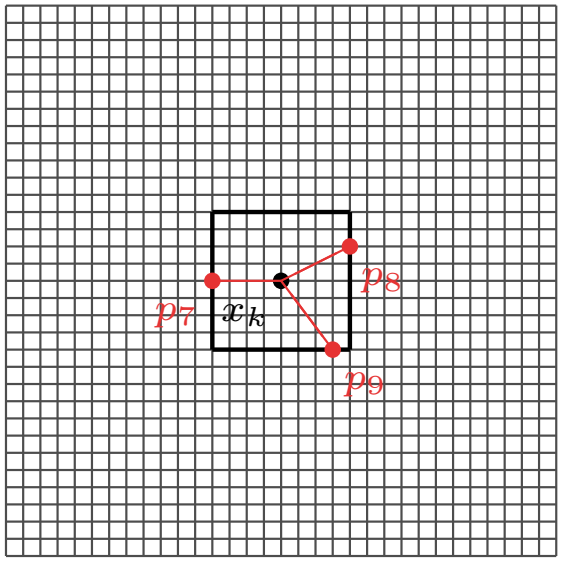


$= \{p_4, p_5, p_6\}$

$\delta^{k+2} = 1/16$
$\Delta^{k+2} = 1/4$



$= \{p_7, p_8, p_9\}$

Discrete variables in MADS – so far

- MADS has been designed for continuous variables.
- Some theory exists for categorical variables [Abramson, 2004].
- So far: Only a patch allows to handle integer variables: Rounding + minimal mesh size of one.

In this talk, we start from scratch and present direct search methods with a natural way of handling discrete variables.

Mesh refinement on $\min(x - 1/3)^2$

Δ^k	x^k
1	0
0.5	0.5
0.25	0.25
0.125	0.375
0.0625	0.3125
0.03125	0.34375
0.015625	0.328125
0.0078125	0.3359375
0.00390625	0.33203125
0.001953125	0.333984375

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alternately

Δ^k	x^k
1	0
0.5	0.5
0.2	0.4
0.1	0.3
0.05	0.35
0.02	0.34
0.01	0.33
0.005	0.335
0.002	0.332
0.001	0.333

Idea: Instead of dividing Δ^k by 2, change it so that

10×10^b refines to 5×10^b

5×10^b refines to 2×10^b

2×10^b refines to 1×10^b

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To get three decimals, one simply sets the granularity to 0.001.
Integer variables are treated by setting the granularity to $\mathcal{G} = 1$.

Poll and mesh size parameter update

- The poll size parameter Δ^k is updated as
  $10 \times 10^b \longleftrightarrow 5 \times 10^b \longleftrightarrow 2 \times 10^b \longleftrightarrow 1 \times 10^b$
- The fine underlying mesh is defined with the mesh size parameter
  $$\delta^k = \begin{cases} 1 & \text{if } \Delta^k \geq 1, \\ \max\{10^{2b}, \mathcal{G}\} & \text{otherwise, i.e. } \Delta^k \in \{1, 2, 5\} \times 10^b. \end{cases}$$
- Example: Granularity of $\mathcal{G} = 0.005$:

δ^k	Δ^k
1	5
1	2
1	1
0.01	0.5
0.01	0.2
0.01	0.1
0.005	0.05
0.005	0.02
0.005	0.01
0.005	0.005 $\leftarrow$ `stop`

[0] Initializations (x_0, $\Delta_0 \in \{1, 2, 5\} \times 10^b$, $\mathcal{G}$ granularity)

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Theory

- MADS analysis relies on "the sequence of trial points are located on some discretization of the space of variables called the mesh".
- By multiplying or dividing Δ^k by a rational number τ, [Torczon, 1997] showed that all trial points from iteration 0 to ℓ were located on a fine underlying mesh. The proof is not trivial and uses the fact that $\tau \in \mathbb{Q}$, and does not work for $\tau \in \mathbb{R}$ (paper won SIAM Outstanding Paper Prize).

Theory

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- By multiplying or dividing Δ^k by a rational number τ, [Torczon, 1997] showed that all trial points from iteration 0 to ℓ were located on a fine underlying mesh.
- With the new mesh, that technical part of the proof becomes:

Consider any trial point t considered from iteration $k = 0$ to ℓ.
If a granularity of $\mathcal{G}_i$ is requested on variable i,

$$t_i \text{ lies on the mesh of granularity } \mathcal{G}_i$$

if no granularity is requested on variable i,

$$t_i \text{ lies on the mesh of granularity } 10^{b_i}$$

with $b_i = \min\{b_i^k : k = 0, 1, \ldots, \ell\}$.

Trust-Region parameters tuning

Find the values of the four parameters $x = (\eta_1, \eta_2, \alpha_1, \alpha_2) \in \mathbb{R}^4_+$ that minimize the overall CPU time to solve 55 `CUTEr` problems.

- The surrogate function s is defined as the time to solve a collection of small-sized problems.
- In 2005, $f(x) \simeq 3h - 4h$ and $s(x) \simeq 1m$. The surrogate is 200 times faster.

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Theorem (Moore's Law is valid)

Year	*CPU* $f(x)$	*CPU* $s(x)$
2005	*13800 s*	*69 s*
2016	*330 s*	*2.3 s*
Ratio:	*42*	*30*

Moore's Law: Speed doubles every 2 years: $2016\text{-}2005 = 11 \Rightarrow$ *5.5 cycles, speedup of* $2^{5.5} \simeq 45$.

Standard $\times 2 \div 2$ *versus* **New** $\{1, 2, 5\} \times 10^b$

0.25	0.75	0.50	2	329.64729
0.25	0.80	0.55	1.55	254.32179
0.25	0.95	0.70	1	NaN
0.35	0.95	0.60	1	NaN
0.15	0.70	0.60	1	NaN
0.30	0.75	0.40	1	NaN
0.40	0.70	0.65	1.55	299.24299
0.05	0.90	0.55	4.70	304.38179
0.25	0.90	0.55	1.55	230.25339
0.25	1.	0.55	1.55	NaN
0.15	0.85	0.60	2.45	332.73279
0.25	1.	0.50	1.10	NaN
0.20	0.95	0.75	1.10	384.09890
0.05	0.90	0.50	1.10	315.53730
0.30	0.85	0.50	1	NaN
0.45	0.70	0.50	5.15	376.13209
0.30	0.90	0.55	1.55	234.08379
0.25	0.90	0.65	1.55	267.51490
0.15	0.90	0.55	2	307.73540
0.25	0.80	0.55	2	344.48180
0.30	0.95	0.55	2.45	284.48099
0.30	0.95	0.45	1	NaN
0.1875	0.9125	0.625	1.775	324.89690
0.2125	0.925	0.5375	1.55	237.63839
0.2625	0.90	0.55	2	297.25700
0.2375	0.8875	0.60	1.55	260.38650
0.225	0.8625	0.5375	1.6625	316.95100
0.3125	0.9375	0.525	1	NaN
0.271875	0.909375	0.559375	1.6625	312.84760
0.25	0.925	0.55	1.55	231.82699
0.26875	0.90	0.546875	1.409375	252.02610
0.246875	0.90	0.571875	1.4375	319.19470
0.234375	0.90	0.5375	1.409375	255.43649
0.25	0.875	0.54375	1.94375	303.55739
0.24375	0.91875	0.54609375	1.59921875	238.37809
0.24609375	0.9109375	0.553125	1.54296875	305.85789
0.26171875	0.903125	0.55234375	1.5640625	269.32650
0.24765625	0.89921875	0.5515625	1.6625	325.77249
0.24921875	0.89609375	0.56171875	1.52890625	252.59420
0.24609375	0.890625	0.53125	1.4515625	313.61869
0.2578125	0.8998046875	0.5529296875	1.5623046875	259.61660
0.251171875	0.9060546875	0.5498046875	1.546484375	252.58099
0.2455078125	0.901171875	0.5515625	1.58515625	295.79539
0.25390625	0.899609375	0.5494140625	1.5939453125	244.50489
0.2515625	0.8998046875	0.5560546875	1.5447265625	251.01300
0.248046875	0.893359375	0.5431640625	1.4796875	265.02609
0.251806640625	0.901220703125	0.55107421875	1.55703125	268.87639
0.248876953125	0.902685546875	0.551123046875	1.553076171875	265.05930
0.252783203125	0.901123046875	0.550439453125	1.54296875	306.20870
0.250830078125	0.900341796875	0.54921875	1.575927734375	315.93150
0.250244140625	0.898876953125	0.552783203125	1.557470703125	261.50459
0.247314453125	0.89697265625	0.546435546875	1.520556640625	262.96580
0.249853515625	0.8998291015625	0.55010986328125	1.55120849609375	239.99700
0.2500732421875	0.90140380859375	0.55042724609375	1.5546142578125	268.07650
0.249755859375	0.9005126953125	0.5485595703125	1.54857177734375	238.33220
0.24979248046875	0.90042724609375	0.55035400390625	1.5370361328125	305.50679
0.25152587890625	0.9000732421875	0.54979248046875	1.547802734375	268.33620
0.24886474609375	0.8975830078125	0.55086669921875	1.56197509765625	264.80709
0.250225830078125	0.900222778320312	0.5490966796875	1.5502197265625	229.52039

0.25	0.75	0.5	2	330.1009
0.25	0.85	0.6	1	NaN
0.15	0.75	0.5	3	323.5899
0.05	0.75	0.5	4	340.3754
0.15	0.75	0.5	2	326.0686
0.05	0.75	0.4	4	330.5359
0.25	0.75	0.5	5	334.6560
0.15	0.85	0.5	3	316.2798
0.15	0.95	0.5	3	303.1888
0.15	1.	0.5	3	NaN
0.15	0.85	0.4	3	318.4759
0.05	1.	0.5	3	NaN
0.15	1.	0.6	3	NaN
0.35	1.	0.5	3	NaN
0.15	1.	0.5	5	NaN
0.05	0.05	0.4	1	NaN
0.15	0.85	0.475	3	313.9554
0.15	1.	0.476	2	NaN
0.15	1.	0.532	4	NaN
0.25	0.95	0.503	3	100286.7
0.15	0.95	0.55	2	303.1791
0.15	0.95	0.6	1	NaN
0.15	0.85	0.45	4	290.1130
0.15	0.75	0.35	6	338.5636
0.05	0.85	0.45	5	310.3020
0.15	0.75	0.45	1	NaN
0.15	1.	0.55	2	NaN
0.05	1.	0.25	3	NaN
0.05	0.55	0.55	3	400.4212
0.35	0.45	0.45	1	391.8666
0.2	0.65	0.45	4	351.8827
0.138	0.85	0.35	4	316.5801
0.151	0.85	0.45	6	291.6727
0.2	0.85	0.45	4	297.7219
0.149	1.	0.45	4	NaN
0.112	0.65	0.55	2	409.2380
0.141	0.85	0.4	5	316.5398
0.15	0.95	0.456	4	283.9472
0.15	1.	0.462	4	NaN
0.2	0.65	0.456	5	346.7783
0.147	0.95	0.556	4	311.6209
0.142	1.	0.456	3	NaN
0.174	1.	0.456	5	NaN
0.1	0.95	0.456	4	280.4177
0.05	0.95	0.456	4	283.5251
0.1	0.85	0.556	4	314.9984
0.1	0.95	0.456	5	281.0185
0.2	0.95	0.456	4	274.8974
0.3	0.95	0.456	4	282.7640
0.2	0.65	0.356	3	344.0217
0.3	0.95	0.456	5	282.3462
0.1	0.95	0.556	4	317.3498
0.2	0.75	0.456	4	329.5962
0.4	0.95	0.556	4	304.3251
0	1.	0.256	3	NaN
0.3	0.85	0.517	4.5	289.2744
0.3	0.95	0.414	3.5	237.7690
0.4	0.95	0.372	3	277.6089

Trust-Region parameter tuning: Results

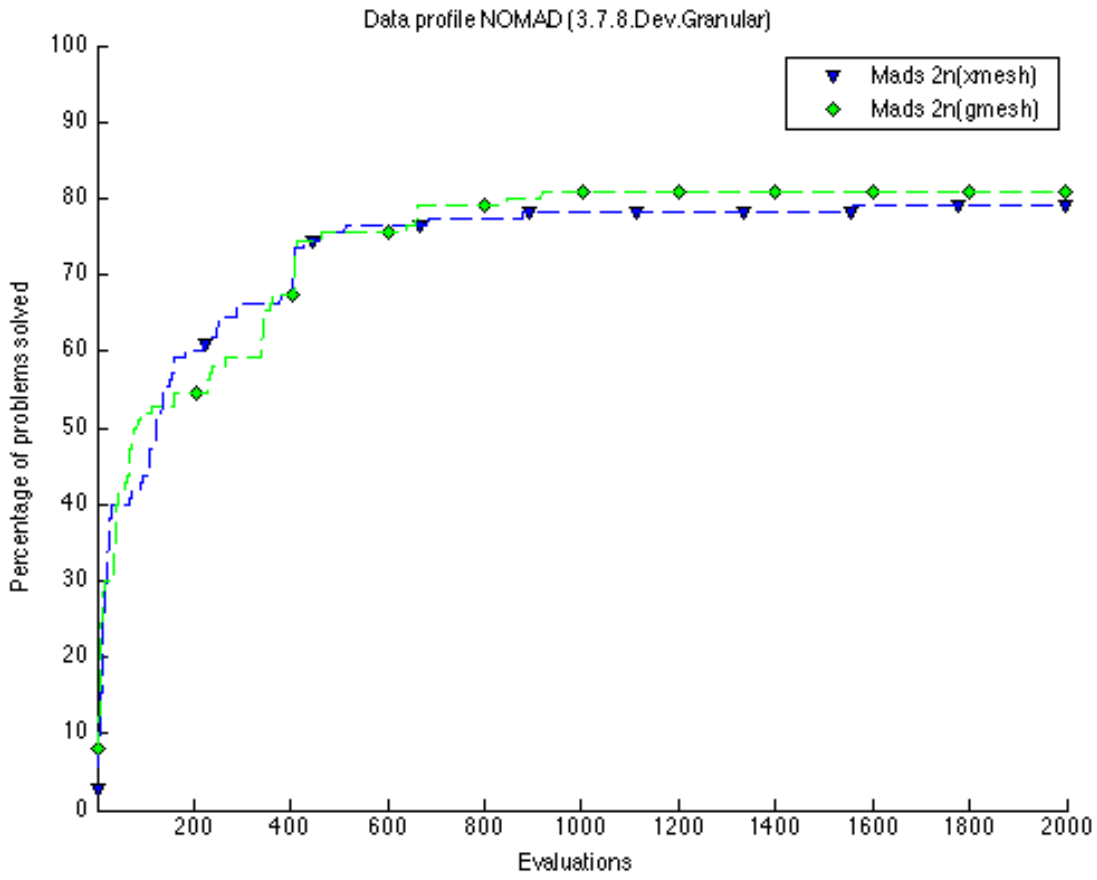
	Standard $\times 2 \div 2$	New $\{1,2,5\} \times 10^b$	Granularity 0.01	Granularity 0.05
CPU:	228 s	207 s	205 s	220 s
η_1	0.2508936274	0.5325498121	0.35	0.35
η_2	0.9008242693	0.9938378210	0.99	0.95
α_1	0.5464425868	0.4817164472	0.50	0.45
α_2	1.5505441560	3.5206901980	2.82	3.20

Observations:

- On this example, the new strategies are preferable.
- Trust-region recommendation for humans:

$$(\eta_1,\ \eta_2,\ \alpha_1,\ \alpha_2) = (0.35,\ 0.99,\ 0.5,\ 2.82).$$

Preliminary results on mixed-integer problems



- Data profiles [Moré and Wild, 2009].
- 11 analytical problems from the literature.
- Continuous+integer variables.
- No constraints.
- 10 random starting points for each problem (110 executions).
- "classic" VS new mesh.

Discussion

- New mesh parameter update rules to control the number of decimals $\{1, 2, 5\} \times 10^b$
  - A native way to handle granularity of variables $\mathcal{G}$.
  - Integer variables are handled by setting $\mathcal{G} = 1$.

Discussion

- New mesh parameter update rules to control the number of decimals $\{1, 2, 5\} \times 10^b$
  - A native way to handle granularity of variables $\mathcal{G}$.
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- Numerical experiments on trust region parameters
  - New parameters reduce CPU time by 35% (versus textbook).
  - New parameters have granularity 0.01 (readable by humans).
  - Experiments on other test problems in progress.

Discussion

- New mesh parameter update rules to control the number of decimals $\{1, 2, 5\} \times 10^b$
  - A native way to handle granularity of variables $\mathcal{G}$.
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- Numerical experiments on trust region parameters
  - New parameters reduce CPU time by 35% (versus textbook).
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  - Experiments on other test problems in progress.
- This will be part of our NOMAD 3.8 software
  - the only additional input from the user is $\mathcal{G}$.
  - ... and it is optional.
  - www.gerad.ca/nomad.

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